

Probabilistic Region Connection Calculus

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Abstract

We present a novel probabilistic model and specification language for spatial relations. Qualitative spatial logics such as RCC are used for representation and reasoning about physical entities. Our probabilistic RCC semantics enables a more expressive representation of spatial relations. We observe that reasoning in this new framework can be hard. We address this difficulty by using a factored representation based on Markov Random Fields.

We provide a logic for representing and reasoning about spatial elements, in the presence of uncertainty. Our framework combines a high-level approach based on qualitative spatial reasoning, that avoids the pitfalls and complexities of pixel-level reasoning, with a probabilistic semantics, able to deal with and quantify uncertainty.

Reasoning about space at the pixel level requires too complex computations and does not capture higher-level properties of objects. As a solution, higher-level qualitative calculi have been introduced, such as Region Connection Calculus, or RCC (Randell, Cui, and Cohn 1992); however, in such calculi there are no shades of gray in representing uncertainty. We take the flexible, high-level approach of qualitative spatial reasoning, RCC-8 in particular, and define probabilistic models.

Using our probabilistic spatial calculus, we are able to answer more accurately questions about the relations between regions: in classic RCC, uncertainty with respect to the base relation that holds between two regions means that some base relations are possible. There is no cue as to which of these relation is more likely. In the worst case, the entire base relation is possible. However, generally, in real world situations, some relations might be more probable than others; using our probabilistic calculus, one can find the probabilities for all the base relations between the two regions and then get, rather than a set of relations, the most probable base relation.

An example of an application for our calculus is recreating a spatial landscape, consisting of all spatial relations that hold between all entities, from a natural language description. The landscape description can be analysed to extract

an initial set of spatial relations as the first, incomplete, landscape, and then the most likely complete image can be recreated using inference in probabilistic RCC. The techniques used here could be extended to other spatial formalisms, that are able to capture other meaningful relations between entities. Reconstructing a spatial landscape from text can be useful to answering deeper understanding queries regarding the text. This kind of queries can nowadays be answered in the context of natural language processing by means of textual entailment (Sammons, Vydiswaran, and Roth 2010). Here, either one uses only lexical cues, which can only lead to a shallow understanding of the text, or one learns to infer deeper, semantic relations implied by the text by training on large corpora of annotated textual entailment pairs. In the latter case, much effort is spent on annotating a corpus and feature engineering. By using qualitative spatial reasoning, one only needs to spend effort in extracting the obvious spatial relations from the text, whereas the deeper understanding queries can be answered by reasoning in the underlying spatial logic.

The paper is structured as follows: first, we present some background notions on RCC. Then, we describe the syntax, semantics, and inference for our calculus. Next, we present the factored representation and inference. Finally, we give an overview of related work and conclude.

Background

Qualitative Spatial Reasoning (Cohn 1997) is a term used for any relational reasoning technique for which the objects are spatial entities.

Region Connection Calculus (RCC), introduced by Randell, Cui and Cohn in 1992 (Randell, Cui, and Cohn 1992), is a qualitative spatial calculus used to reason about the relations between regions. The distinction between base relations is made based on either connectedness or the mereological 'part of' relation. The two definitions are equivalent, as the two relations can be defined by means of each other. Given the possible distinctions and additional information considered (e.g., whether the region borders are taken into account or not), the space of possible relations is broken into a set of jointly exhaustive and pairwise disjoint, or JEPD, base relations.

The RCC-8 base relations are: disconnected (DC), externally connected (EC), partially overlap (PO), tangen-

tial proper part and its inverse (TPP, TPPI), non-tangential proper part and its inverse (NTPP, NTPPI), and equivalent (EQ).

For RCC-5, the border information is not considered, and consequently B is $\{DC, O, EQ, PP, PPI\}$.

Probabilistic RCC

We are given a set of regions in a topology, a set of region names or region constants, and a set of spatial constraints on them expressed as a formula. We want to answer queries regarding the probability of certain relations to hold between certain pairs of regions.

Syntax

The signature of probabilistic RCC is a first-order logic signature of a particular form, containing: a set of constants \mathcal{C} (the region names); and a set of arity 2 relations \mathcal{B} (the base relations).

Two probabilistic RCC-8 (RCC-5) signatures may differ from each other on their set of constants, so:

Definition 1 A PRCC signature is a set of region constants.

Henceforth we will refer to RCC-8 only; the results can easily be applied to RCC-5.

Definition 2 The basic sentences of probabilistic RCC-8 are defined inductively as follows:

- atoms are of the form $r(a, b)$, where $a, b \in \mathcal{C}$ and $r \in \mathcal{B}$;
- if ϕ and ψ are basic RCC-8 sentences, then $\phi \vee \psi$ and $\phi \wedge \psi$ are also basic RCC-8 sentences;
- if ϕ is a basic RCC-8 sentence, then $\neg\phi$ is also a basic RCC-8 sentence

A basic sentence encodes the constraints for the problem and is just a ground FOL sentence. The queries are on the probability of a relation to hold between two regions. This relation may be either a base relation (EC) or a general relation (a disjunction of base relations). For example, 'part-of' is the disjunction of EQ, TPP, and NTPP.

One property of PRCC sentences, that stems from JEPD-ness, namely the fact that the negation of a literal can be rewritten as a positive disjunction, is the following:

Property 1 Any basic sentence of probabilistic RCC-8 can be written as a positive sentence

A conditional query-type sentence expresses the probability of a relation given a basic type sentence: this is the kind of sentence that generally encodes a full problem. The semantics of these sentences is defined using the semantics of non-conditional query-type sentences.

In the following, α is the probability we are looking for: $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$ has the intuitive meaning that the probability that $r(a, b)$ holds is α .

Definition 3 If $0 \leq \alpha \leq 1$, $a, b \in \mathcal{C}$, $\mathcal{B}_q \subset \mathcal{B}$ and ϕ is a basic sentence, then:

- $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$ is a non-conditional query-type sentence or a query-type atom;
- $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ is a conditional query-type sentence.

Definition 4 A probabilistic RCC-8 sentence is either a basic sentence or a query-type sentence.

Semantics

A model of a PRCC signature will specify the topology, a subset of this topology (the 'working' regions), the set of interpretations of region constants in the 'working' region set and a probability distribution on these interpretations.

Let T be a topology on some universe U and let $X \in \mathcal{R}$ be a closed set in T . In the following, let $Int(X)$ be the interior of X and $\Gamma(X) = X - Int(X)$ be the border of X .

Definition 5 Given an RCC-8 signature \mathcal{C} , a model M is a structure of the form $M = (U, T, \mathcal{R}, W, P)$, where:

- U is a (possibly infinite) universe of points;
 - T is a topology on U ; the closed regular sets in T are regions;
 - $\mathcal{R} \subset T$ is a finite set of regions;
 - $W = \{(U_w, w) \mid w : \mathcal{C} \uplus \mathcal{B} \rightarrow U_w \uplus (U_w \times U_w)\}$ is a set of possible worlds, where for each possible world w :
 - $U_w = \mathcal{R}$ is the world universe;
 - $w|_{\mathcal{C}} : \mathcal{C} \rightarrow U_w$ is an interpretation of constant symbols as regions;
 - $w|_{\mathcal{B}} : \mathcal{B} \rightarrow U_w \times U_w$ is an interpretation of base relation symbols
 - $P : W \rightarrow [0, 1]$ (with $\sum_{w \in W} P(w) = 1$) is a probability distribution over the set of interpretations.
- and the interpretation of base relation symbols $w|_{\mathcal{B}}$ is such that $\forall X, Y \in U_w$:
- $w(DC)(X, Y)$ iff $X \cap Y = \emptyset$;
 - $w(EC)(X, Y)$ iff $Int(X) \cap Int(Y) = \emptyset$ and $X \cap Y \neq \emptyset$;
 - $w(PO)(X, Y)$ iff $Int(X) \cap Int(Y) \neq \emptyset$ and $X \not\subseteq Y$ and $Y \not\subseteq X$;
 - $w(EQ)(X, Y)$ iff $X = Y$;
 - $w(TPP)(X, Y)$ iff $X \subsetneq Y$ and $X \not\subseteq Int(Y)$;
 - $w(TPPI)(X, Y)$ iff $w(TPP)(Y, X)$;
 - $w(NTPP)(X, Y)$ iff $X \subseteq Int(Y)$;
 - $w(NTPPI)(X, Y)$ iff $w(NTPP)(Y, X)$.

These properties also ensure that the set $w(\mathcal{B})$ forms a partition over $U_w \times U_w$, or in other words the relations in $w(\mathcal{B})$ are JEPD.

In what follows, we will assume the topological space fixed. The interpretation of base relations in this space will be the same for all models so we will omit both of these elements. Moreover, for all models we will have the set of interpretations to be the entire set of functions from \mathcal{C} to \mathcal{R} , so W will be completely defined by \mathcal{R} and can be omitted as well (restrictions to a subset of \mathbb{I} can be made by forcing the probability of the missing interpretation functions to 0).

For basic sentences, sentence satisfaction is defined for every possible world, inductively on the structure of the sentence, as in any fragment of FOL. A sentence is satisfied if it is satisfied in every world that has a non-zero probability.

Definition 6 Given model $M = (\mathcal{R}, W, P)$, the satisfaction of a basic formula in a possible world $w \in W$ is defined inductively as:

- $w \models r(a, b)$ iff $(w(a), w(b)) \in w(r)$;
- $w \models \phi \wedge \psi$ iff $w \models \phi$ and $w \models \psi$;
- $w \models \neg\phi$ iff $w \not\models \phi$;
- $w \models \phi \vee \psi$ iff $w \models \neg(\neg\phi \wedge \neg\psi)$;

We say a model $M = (\mathcal{R}, W, P)$ satisfies a basic formula ϕ and write $M \models \phi$ iff $w \models \phi$ for all $w \in W$ with $P(w) > 0$.

Next, we will show how to answer queries, given a model and a set of constraints. The intuition is that, when we are presented with a new piece of information about the world, we constrain our model of the world so as to discard all interpretations that are not consistent with the new piece of information. The model we end up with is what we will call the restriction of a model via a basic-type sentence. Restricting the model via a sentence lowers to 0 the probabilities of all the interpretations that do not satisfy the sentence, and scales the other probabilities such that they still sum to 1.

Definition 7 Let ϕ be a basic formula and $M = (\mathcal{R}, W, P)$ a probabilistic RCC-8 model; then we can define the restriction of M via ϕ as $M|_\phi = (\mathcal{R}, W, P|_\phi)$, where:

- $P|_\phi(w) = P(w) \cdot \frac{1}{Z(\phi)}$ if $w \models \phi$;
- $P|_\phi(w) = 0$ if $w \not\models \phi$

and $Z(\phi) = \sum_{w \models \phi} P(w)$ is the normalization constant.

Thus $M|_\phi$ is intuitively the largest submodel of M that satisfies ϕ .

In order to answer the query given a set of constraints, we restrict the model to satisfy the set of constraints, and then we sum the probabilities of the interpretations that satisfy the query. So, the satisfaction of a query-type sentence by a model M is defined as follows:

Definition 8 Given model $M = (\mathcal{R}, W, P)$, basic sentence ϕ , $a, b \in \mathcal{C}$, $\mathcal{B}_q \subset \mathcal{B}$ and $0 \leq \alpha \leq 1$, the satisfaction of query-type sentence $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ is defined as:

- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$ iff $\sum_{w \models \bigvee_{r \in \mathcal{B}_q} r(a, b)} P(w) = \alpha$;
- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ iff $M|_\phi \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$.

It is worth noting that we are really not interested in what exactly the interpretations of constant symbols in a possible world look like, but in their relative position. So we can restrict our attention to equivalence classes of possible worlds, under the equivalence relation \simeq given by the set of base RCC relations that hold in these worlds:

$$w_1 \simeq w_2 \quad \text{iff} \quad \text{for each pair } a, b \in \mathcal{C} \text{ and for each } r \in \mathcal{B} \\ w_1 \models r(a, b) \Leftrightarrow w_2 \models r(a, b) \quad (1)$$

Inference in Probabilistic RCC

Using definitions 6 and 8, we can derive the following alternative condition for the satisfaction of conditional query-type sentences - $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$:

$$\alpha = \frac{\sum_{w \models \phi \text{ and } w \models \bigvee_{r \in \mathcal{B}_q} r(a, b)} P(w)}{\sum_{w \models \phi} P(w)} \quad (2)$$

It is straightforward to implement an algorithm that finds α using this formula. If N is the size of ϕ , R is the number of regions and C is the number of constant symbols in the signature, this algorithm would require $O(R^{C+1})$ space and $O(N \cdot R^C)$ time.

Notice that this algorithm requires us to know P , the probability distribution over possible worlds. If we don't know it, the proper probability distribution to use is the one with the maximum entropy, according to the principle of maximum entropy. The set of possible worlds being a discrete and finite domain, the maximum entropy probability distribution is the uniform probability distribution.

Using this observation and equation (2), we can compute α in $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ as:

$$\alpha = \frac{|\{w \in W \mid w \models \phi \text{ and } w \models \bigvee_{r \in \mathcal{B}_q} r(a, b)\}|}{|\{w \in W \mid w \models \phi\}|} \quad (3)$$

Factored Representation of PRCC

Given a signature \mathcal{C} and model $M = (\mathcal{R}, W, P)$, for each pair of distinct constant symbols $a, b \in \mathcal{C}$, let $X_B^{a,b}$ be the random variable encoding the base relation that holds between the regions named by a and b . Then, the probability distribution P over possible worlds induces a joint probability P_B distribution over $\{X_B^{a,b}\}_{a,b \in \mathcal{C}, a \neq b}$:

$$P_B(X_B^{p_1} = r_1, \dots, X_B^{p_N} = r_N) = \sum_{w \models \bigwedge_{1 \leq i \leq N} r_i(p_i)} P(w) \quad (4)$$

where $N = \binom{\mathcal{C}}{2}$, $\{p_1, \dots, p_N\} = \{\{a, b\} \in \mathcal{C} \mid a \neq b\}$ and $r_i \in \mathcal{B}$ for $1 \leq i \leq N$.

If we consider the model consisting of equivalence classes of possible worlds, we can recover the probability distribution over these equivalence classes from the joint probability P_B , as:

$$P([w]_{r_1(p_1), \dots, r_N(p_N)}) = P_B(X_B^{p_1} = r_1, \dots, X_B^{p_N} = r_N) \quad (5)$$

where $[w]_{r_1(p_1), \dots, r_N(p_N)} = \{w \in W \mid w \models r_1(p_1) \wedge \dots \wedge r_N(p_N)\}$.

Therefore, reasoning in probabilistic RCC can be reduced to reasoning with such joint probability distributions:

Theorem 1 Given model $M = (\mathcal{R}, W, P)$, basic sentence ϕ , expressed as a conjunction of atoms, $a, b \in \mathcal{C}$, $r \in \mathcal{B}$ and $0 \leq \alpha \leq 1$, the satisfaction of query-type sentence $p_\alpha(r(a, b) \mid \phi)$ can be computed as follows:

- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$ iff $P_B(X_B^{a,b} = r) = \alpha$;
- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ iff $P_B(X_B^{a,b} = r \mid \phi) = \alpha$;

Furthermore, since for every $X_B^{a,b}$, the events $X_B^{a,b} = r$ and $X_B^{a,b} = r'$ are disjoint for every $r \neq r' \in \mathcal{B}$:

Corollary 1 Given model $M = (\mathcal{R}, W, P)$, basic sentence ϕ , expressed as a conjunction of atoms, $a, b \in \mathcal{C}$, $\mathcal{B}_q \subset \mathcal{B}$ and $0 \leq \alpha \leq 1$, the satisfaction of query-type sentence $p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ can be computed as follows:

- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b))$ iff $\sum_{r \in \mathcal{B}_q} P_B(X_B^{a,b} = r) = \alpha$;
- $M \models p_\alpha(\bigvee_{r \in \mathcal{B}_q} r(a, b) \mid \phi)$ iff $\sum_{r \in \mathcal{B}_q} P_B(X_B^{a,b} = r \mid \phi) = \alpha$.

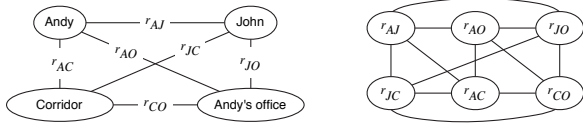


Figure 1: Compact representation of a model

Markov Random Fields

We can represent the joint probability distribution P_B as a Markov Random Field (MRF), using the observation that the base relation that holds between two regions named a and b is independent of any other relation that holds in the world, given the relations that hold between region named a and any other region, and the relations that hold between any other region and region named b , and these relations' duals:

$$I(X_B^{a,b}, X_B^{c_i, c_j} \mid X_B^{a,c} \cup X_B^{b,c} \cup X_B^{c,a} \cup X_B^{c,b}) \quad (6)$$

where $X_B^{a,c} = \{X_B^{a,c} \mid c \in \mathcal{C}\}$ and $X_B^{c,a} = \{X_B^{c,a} \mid c \in \mathcal{C}\}$.

This observation does not hold for all PRCC models, and we will only be able to use this compact representation for those models that do have this property. Intuitively, this is the case if we don't have any prior knowledge of the space of regions \mathcal{R} , and indeed in this case, using the naive algorithm described in the previous section is infeasible.

The MRF representation for unstructured models is illustrated on a simple example in Figure 1. Notice that if (6) holds, then we only have edges between nodes that share a symbol.

Lemma 1 *If the assumption 6 holds for a model M , then in the MRF representation of M , there is an edge between nodes X_B^p and X_B^q if pairs p, q share a constant symbol, i.e. $|p \cap q| = 1$.*

Theorem 2 *Let \mathcal{C} be a probabilistic RCC signature and let $P_B(X_B^{p_1} = r_1, \dots, X_B^{p_N} = r_N)$ be the probability distribution over the base relations that hold between the interpretations of every two constant regions, given an unstructured model M . Then, in the MRF representation of P_B , the largest clique has size $C - 1$.*

Intuitively, every node $X_B^{a,b}$ ($a \neq b \in \mathcal{C}$) in the MRF representation is connected to two cliques of size $C - 1$: one containing all the pairs that share symbol a , and one containing all the pairs that share symbol b . Other cliques that appear in the MRF are triangles representing the relations that hold between any three regions. The interactions represented by those latter cliques stem from the constraints imposed by RCC relation composition.

Let $\bar{X}_B^{a,?} = (X_B^{a,b})_{b \neq a, b \in \mathcal{C}}$ be the tuple containing the nodes in the clique sharing symbol a , and let $X_B^{a,b,c} = X_B^{a,b}, X_B^{b,c}, X_B^{c,a}$. For an unstructured model, one can have any combination of base relations between a region and all the other regions, i.e., we can assume $\phi(\bar{X}_B^{a,?})$ a constant and therefore ignore it in the factorization. Therefore the probability distribution can be written as:

$$P_B(X_B^{p_1} = r_1, \dots, X_B^{p_N} = r_N) = \frac{1}{Z_B} \prod \phi_{a,b,c}(X_B^{a,b,c}) \quad (7)$$

Inference in the Factored Models

In the following we will assume we know the factors $\phi_{a,b,c}(X_B^{a,b}, X_B^{b,c}, X_B^{c,a})$ in the joint probability distribution. We can infer the probability α of $\forall r \in \mathcal{B}_q r(a, b)$ as the sum of probabilities of each $r(a, b)$, given an evidence $\phi = r_1(a_1, b_1) \wedge \dots \wedge r_k(a_k, b_k)$:

$$\alpha = \frac{\sum_{r \in \mathcal{B}_q} P(r(a, b), r_1(a_1, b_1), \dots, r_k(a_k, b_k))}{\sum_{r \in \mathcal{B}} P(r(a, b), r_1(a_1, b_1), \dots, r_k(a_k, b_k))} \quad (8)$$

If we further assume

$$\phi_{a,b,c}(X_B^{a,b}, X_B^{b,c}, X_B^{c,a}) = w_{a,b,c} f_{a,b,c}(X_B^{a,b}, X_B^{b,c}, X_B^{c,a})$$

where the value of the feature $f_{a,b,c}$ is 1 if the configuration specified by the relations between a, b and c is possible and 0 otherwise, we can use any MRF learning algorithm to infer the set of weights $\{w_{a,b,c}\}_{a,b,c}$. For the current work we assume that all factors are known, or all weights are 1.

Theorem 3 *Variable Elimination for the factored model of PRCC has a time complexity of $O(2^{C^3})$.*

Related Work

Another way to do probabilistic reasoning in RCC is to use the language of Markov Logic Networks (Richardson and Domingos 2006). This amounts to representing PRCC as an MLN built from an axiomatization (Randell, Cui, and Cohn 1992) of classic RCC. What we do here is to give probabilistic RCC an individuality of its own, with its own well-defined syntax and semantics. We also encode the PRCC models directly as MRFs, taking advantage of the particular independence assumptions that stem from the spatial domain.

(Cohn and Gotts 1995) address the problem of reasoning in the presence of vague topological information, in the case where the regions have vague boundaries. They introduce the 'egg-yolk' representation, where each region is divided into its crisp, certain subregion (the 'yolk') and a surrounding vague part (the 'white'). The intuition is that the actual region lies anywhere within the borders of the 'white' and necessarily covers the 'yolk'. In this work there is no quantification for the degree of uncertainty.

(Schockaert and De Cock 2007) and (Schockaert et al. 2008) also deal with vague regions and add quantifiable uncertainty. Rather than work in a probabilistic setting, as in our approach, or divide each region, as in the previous approach, they develop a framework based on fuzzy logic. They take the 'connected' relation to mean the degree to which regions are connected, not a crisp truth value as in the classical RCC. With this, they redefine the entire set of base relations of RCC and subsequently the RCC framework. In contrast, we keep the classic logical framework of RCC and give it a probabilistic semantics.

In order to deal with uncertainty regarding regions, (Bitner and Stell 2001) represent approximate regions by relating them to a frame of reference consisting of a set of unit regions. The definition of approximation makes qualitative distinctions based on the coverage of those unit regions. They then define an approximate region as a set of

regions with the same approximation. With this definition, they rewrite the RCC framework to work with approximate regions.

All of these lines of work look at dealing with or quantifying vagueness rather than quantifying the likelihood of relations.

Probabilistic logic programming, or PLP (Lukasiewicz 1998), resembles our work mainly in the way they define the satisfaction of probabilistic sentences. One major difference is that, in PLP, each probabilistic formula (representing the probability of a conditional event) is assigned a probability *interval* - we are reasoning over single probabilities, not probability intervals. Another important difference is that any sentence in PLP is a probabilistic sentence; in our case, only the queries are probabilistic, whereas the knowledge base consists only of sentences expressed in classic RCC.

A maximum entropy semantics has also been defined for PLP (Yue, Liu, and Hunter 2008); that definition is based on the notion of degree of satisfaction. Since we do not use probabilistic sentences in our knowledge base, our maximum entropy model is much more simple.

(Kontchakov, Pratt-Hartmann, and Zakharyashev 2010) have proved the sensitivity to the underlying topological space of the complexity of reasoning in a superset of RCC-8, enriched with a unary connectedness predicate and Boolean functions over regions. They also prove NP-completeness of satisfiability for the calculus enriched with connectedness only, as well as EXPTIME-hardness of the full superset. Further results (Kontchakov et al. 2011) prove reasoning in the 2D Euclidean space RE-hard for the case when Boolean functions are allowed over regions. We do not make assumptions on the underlying topological space, and we do not talk about Boolean functions over regions.

Conclusions and Further Work

We showed PRCC syntax, semantics, and compact representation using Markov Random Fields to model the joint probability distribution over spatial relations.

One problem we don't address is how to handle disjunctive evidence. One way to look at this, is that, writing the evidence in disjunctive normal form, every clause serves as evidence for a possible abstract image of the world. We would then want to combine the probabilities of base relations that result from each of these possible images. One could take an optimistic approach and use the maximum of these probabilities, for every query, but this does not accurately represent the probability distribution encoded by the model. We will explore ways to look at disjunctive evidence in future work.

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