Reasoning about Truthfulness of Agents Using Answer Set Programming

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Abstract
We propose a declarative framework for representing and reasoning about truthfulness of agents using answer set programming. We show how statements by agents can be evaluated against a set of observations over time equipped with our knowledge about the actions of the agents and the normal behavior of agents. We illustrate the framework using examples and discuss possible extensions that need to be considered.

Introduction
The recent advancements of AI enable the development of agents that can replace human in many tasks. Increasingly, organizations are using web-bots for interacting with clients in various capacities, e.g., in providing information about the company or making offers. It is reasonable to believe that this trend will be continued as long as the Internet exists. Unfortunately, not every business on the Internet is honest as one would hope. Stories about businesses that cheat people of goods or services or wrongly advertise their services are not uncommon. This leads to the development of businesses that allow people to rate companies (e.g., expedia.com or Angie’s List) or defend the reputation of a company or an entity (e.g., reputation.com).

In this paper we are interested in reasoning about the truthfulness of agents. We start with an optimistic assumption that agents are truthful unless otherwise proven. We will judge agents by what they do or what we observe rather than by what they say. Our observations are made at different time instances along a linear time line. We assume that whatever observed is true at the time it is observed and will stay true until additional information indicates otherwise. Furthermore, we will need to judge agents even when we do not have complete information about them. This means that reasoning about the truthfulness of agents is a non-monotonic task that is often done under incomplete information. To illustrate these issues, let us consider the following situations.

1. John said that he does not have money.
   We observe that John does not attend college.
   Given the above observation, we would likely conclude that John is being honest in his statement, i.e., we trust that John does not have money which does not allow him to attend a college.

2. We observe that John has a Ferrari.
   A common person would normally be unable to afford a luxury car. At this point, we will likely withdraw our conclusion about John having no money and classify John as a liar.

3. We learn that John inherited the Ferrari from his parents.
   Inheritance is a possible way to acquire something and thus we might need to reconsider our position about John being a liar. Yet, if someone inherits a Ferrari from his parents then commonsense dictates that the person comes from an affluent household and thus should have money (enough to attend a college). All this reasoning implies that we should not trust John even after this new piece of information.

4. We learn that John is a gambler.
   A gambler is usually penny-less. Luxury items he may own are likely obtained by gambling rather than purchased. Hence, we might be inclined to trust that John does indeed have no money.

   This sequence of observations and the associated reasoning processes are rather simple, but they illustrate clearly the way humans reason about the truthfulness of statements and make judgements on them. The observation that John does not attend college is used to deduce that John does not have (enough) money to attend college. This makes John’s statement truthful. However, observations such as “John has a Ferrari” or “John inherited a Ferrari from his family” would cause us to withdraw our conclusion about John being poor, i.e., John’s statement is now untruthful. Yet, an additional observation that John is a gambler will force us to withdraw our latest conclusion and conclude again that his statement is truthful. Understandably, this process can continue forever and our opinion about John’s statement might or might not change after each round of new observations.

In this paper, we propose a framework for representing and reasoning about the truthfulness of agents using Answer Set Programming, a knowledge representation language useful for commonsense reasoning in presence of incomplete information (Baral 2003).
Background: Answer Set Programming

In this section, we will review the basic definitions of Answer Set Programming (ASP) (Baral 2003), i.e., the language of logic programs under the answer set semantics (Gelfond and Lifschitz 1990). A logic program $\Pi$ is a set of rules of the form

$$c_1 \mid \ldots \mid c_k \leftarrow a_1, \ldots, a_m, \neg a_{m+1}, \ldots, \neg a_n$$  \hspace{1cm} (1)

where $0 \leq m \leq n$, $0 \leq k$, each $a_i$ or $c_j$ is a literal of a propositional language\(^1\) and $\neg$ represents negation-as-failure. A negation as failure literal (or naf-literal) is of the form $\neg a$ where $a$ is a literal. For a rule of the form (1), the left and right hand sides of the rule are called the head and the body, respectively. Both the head and the body can be empty. When the head is empty, the rule is called a constraint. When the body is empty, the rule is called a fact.

For a rule $r$ of form (1), $H(r)$ and $B(r)$ denote the left and right hand side of $\leftarrow$, respectively; $\text{head}(r)$ denotes $\{c_1, \ldots, c_k\}$; and $\text{pos}(r)$ and $\text{neg}(r)$ denote $\{a_1, \ldots, a_m\}$ and $\{a_{m+1}, \ldots, a_n\}$. For a program $\Pi$, lit$(\Pi)$ denotes the set of literals occurring in $\Pi$.

Consider a set of ground literals $X$. $X$ is consistent if there exists no atom $a$ such that both $a$ and $\neg a$ belong to $X$. The body of a rule $r$ of the form (1) is satisfied by $X$ if $\text{neg}(r) \cap X = \emptyset$ and $\text{pos}(r) \subseteq X$. A rule of form (1) with nonempty head is satisfied by $X$ if either its body is not satisfied by $X$ or $\text{head}(r) \cap X \neq \emptyset$. A constraint is satisfied by $X$ if its body is not satisfied by $X$.

For a consistent set of ground literals $S$ and a program $\Pi$, the reduct of $\Pi$ w.r.t. $S$, denoted by $\Pi^S$, is the program obtained from the set of all ground instances of $\Pi$ by deleting (i) each rule that has a naf-literal $\neg a$ in its body with $a \in S$, and (ii) all naf-literals in the bodies of the remaining rules.

$S$ is an answer set (or a stable model) of $\Pi$ (Gelfond and Lifschitz 1990) if it satisfies the following conditions: (i) If $\Pi$ does not contain any naf-literal (i.e. $m = n$ in every rule of $\Pi$) then $S$ is a minimal consistent set of literals that satisfies all the rules in $\Pi$; and (ii) If $\Pi$ does contain some naf-literal ($m < n$ in some rules of $\Pi$), then $S$ is an answer set of $\Pi$ if $S$ is the answer set of $\Pi^S$. Note that $\Pi^S$ does not contain naf-literals; hence its answer set is defined in the first item. A program $\Pi$ is said to be consistent if it has a consistent answer set. Otherwise, it is inconsistent.

Answer sets of logic programs can be computed using answer set solvers such as CLASP (Gebser, Kaufmann, and Schaub 2009).

An Illustration

In this section, we encode the reasoning process that is used in the introduction to evaluate John’s statements. We assume a propositional language that encodes properties of interest such as $\text{has\_money}$ (John has money), $\text{in\_college}$ (John is in college), etc. with obvious meaning\(^2\). In addition, we will use special predicates of the form $\text{statement}(x, t)$ to denote that “the statement $x$ is made at time step $t$,” $\text{observed}(x, t)$ to denote that “we observe $x$ is true at the time step $t$,” and $h(x, t)$ to indicate that “$x$ is (believed to be) true at the time step $t.” For simplicity, the representation of time is based on discrete steps with negligible duration.

John’s statement that he does not have money can be encoded by the following formula:\(^3\)

$$\text{statement}(\neg \text{has\_money}, 0).$$  \hspace{1cm} (2)

Our observation that John does not attend college is encoded as follows.

$$\text{observed}(\neg \text{in\_college}, 0).$$  \hspace{1cm} (3)

The common knowledge that attending college normally requires money (i.e., normally, if a person does not attend college, then he/she does not have money) can be encoded as follows.

$$h(\neg \text{has\_money}, T) \leftarrow h(\neg \text{in\_college}, T), \text{not}\ ab_1(T).$$  \hspace{1cm} (4)

In the above rule, $\text{ab}_1(T)$ is used to characterize possible exceptions to the rule. For example,不合格 individuals or people who are not interested in a college degree are exceptions to this rule.

Let $\pi_1$ be a program consisting of the commonsense rule (4), the observation (3), and the rules

$$h(l, N) \leftarrow \text{observed}(l, T), N \geq T, \text{not}\ ab_1(l, T, N).$$ \hspace{1cm} (5)
$$\neg ab_1(l, T, T) \leftarrow \text{observed}(l, T)$$ \hspace{1cm} (6)
$$ab_1(l, T, N) \leftarrow h(l, T, N).$$ \hspace{1cm} (7)

that we will refer to as rules for reasoning about observations. $l$ denotes the complement of $l$, i.e., $\overline{\pi} = \neg a$ and $\neg \overline{a} = a$ for an atom $a$. Rule (5) says that an observation about $l$ (a literal) is true from the time it is observed unless otherwise stated. Rule (6) indicates that an observation must be true at the time it is made. Finally, (7) says that newer information will override older one. It is easy to see that the program $\pi_1$ has a unique answer set $S_1$ consisting of (3) and

$$\{h(\neg \text{in\_college}, 0), h(\neg \text{has\_money}, 0), \neg ab_1(\text{has\_money}, 0, 0)\}$$

Furthermore, $h(\neg \text{has\_money}, 0)$ indicates that John does not have money (at step 0), which means that the literal occurs in the statement (2) is “satisfied” by $S_1$.

Let us consider the second observation. We observe that John has a Ferrari. This is expressed as follows.

$$\text{observed}(\text{has\_ferrari}, 1).$$  \hspace{1cm} (8)

\(^1\)Rules with variables are viewed as a shorthand for the set of their ground instances.

\(^2\)For simplicity, we will not attach the subject (John) to the literals as our discussion focuses on him. The programs and discussions can easily be extended to consider several subjects simultaneously.

\(^3\)We abuse notation slightly and use $\neg$ in conjunction with reified properties.
A person with a Ferrari usually has money. This is encoded as follows4.

\[ h(\text{has money}, T) \leftarrow h(\text{has ferrari}, T), \text{not } ab_{2}(T). \quad (9) \]

It is easy to see that rules (5)–(6) will allow us to conclude that \( h(\text{has money}, 1) \) is true in every answer set of the program \( \pi_{2} \) that consists of \( \pi_{1} \) and (8)–(9). In fact, it can be shown that the unique answer set \( S_{2} \) of \( \pi_{2} \) is the union of \( S_{1}, (8) \), and the two literals

\[ \{ h(\text{has money}, 1), h(\text{has ferrari}, 1), \}
\[ \{ \text{not } ab_{2}(\text{has money}, 0, 1) \} \]

which indicates that, at step 1, we believe that John does have money, i.e., John’s statement is not truthful.

Let us consider the next observation that John inherited the Ferrari.

\[ \text{observed}(\text{has ferrari from parent}, 2). \quad (10) \]

Our commonsense knowledge stating that if a person inherits a Ferrari then that person would come from an affluent family and thus would have money is encoded—in a simplified manner—as follows.

\[ h(\text{has money}, T) \leftarrow h(\text{has ferrari from parent}, T), \text{not } ab_{3}(T). \quad (11) \]

Let \( \pi_{3} \) be the program that consists of \( \pi_{2} \) and (10)–(11). Again, we can easily check that this program has a unique answer set \( S_{3} \) that consists of \( S_{2}, (10) \), and the two literals

\[ \{ h(\text{has money}, 2), h(\text{has ferrari from parent}, 2) \}
\[ \{ \text{not } ab_{3}(\text{has money}, 0, 2) \} \]

which indicates that John has money and thus is not truthful.

Let us now consider the last observation that John is a gambler. This is expressed by the fact

\[ \text{observed}(\text{is gambler}, 3). \quad (12) \]

and our knowledge about gamblers is expressed by

\[ h(\text{has money}, T) \leftarrow h(\text{is gambler}, T), \text{not } ab_{4}(T). \quad (13) \]

A gambler would constitute an exception to the rule (9) as well as to the rule (11). This means that we also need the following rules

\[ ab_{2}(T) \leftarrow h(\text{is gambler}, T). \quad (14) \]

\[ ab_{3}(T) \leftarrow h(\text{is gambler}, T). \quad (15) \]

Let \( \pi_{4} \) be the program consisting of \( \pi_{3} \) and the rules (12)—(15). We can easily check that it has one answer set \( S_{4} \) that consists of \( S_{3}, (12) \), and the literals

\[ \{ ab_{2}(3), ab_{3}(3), h(\text{is gambler}, 3), h(\text{has money}, 3) \}
\[ \{ \text{not } ab_{4}(\text{has money}, 1, 3), \text{not } ab_{4}(\text{has money}, 2, 3) \} \]

This answer set allows us to conclude that, at step 3, we believe that John's initial statement is indeed true.

\[ \text{Our Framework} \]

The previous section provides a glimpse on how answer set programming can be used in reasoning about the truthfulness of agents. We will now present a framework for representing and reasoning about the truthfulness of (statements made by) agents. We assume that

- agents’ actions and their effects are fully observable, and actions have preconditions (e.g., the action of buying a car requires that the agent has money and its execution will result in the agent owning a car);
- we can observe over time the occurrence of the agents’ actions and the properties of the world (e.g., observe that John buys a car; John is a gambler);
- we have a repertoire of commonsense knowledge about normal behaviors (e.g., a person buying a luxury car normally has money; a gambler usually has no money).

This leads us to define a knowledge base about an agent as follows.

**Definition 1** A knowledge base (about an agent) over a propositional language \( \mathcal{L} \) is a triple \( \langle \Pi_{A}, \Pi_{K}, \Pi_{O} \rangle \) where

- \( \Pi_{A} \) is a program encoding the agents’ actions and their effects by defining the truth value of literals of the form \( h(l, t) \) where \( l \) is a (reified) literal of the language \( \mathcal{L} \) and \( t \) denotes a time step.
- \( \Pi_{K} \) is a program encoding commonsense knowledge that contains rules defining \( h(l, t) \) and \( \Pi_{O} \) is a collection of atoms of the form \( \text{observed}(l, t) \) or \( \text{occurred}(a, t) \) where \( l \) is a literal of \( \mathcal{L} \), \( a \) is an action defined in \( \Pi_{A} \), and \( t \) is a time step.

Observe that there is extensive literature on reasoning about actions and change and commonsense reasoning using logic programming (e.g., (Baral and Gelfond 1994; Gelfond and Lifschitz 1992; Son et al. 2006)). As such, the two components \( \Pi_{A} \) and \( \Pi_{K} \) can be adapted from previous research for the purpose of this paper in a straightforward manner.

For instance, \( \Pi_{A} \) can follow the proposal given in the seminal paper (Gelfond and Lifschitz 1992); \( \Pi_{K} \) can contain the rules (9), (11), (13), and (14)–(15), etc. discussed in the previous section.

Next, we define the semantics of knowledge bases. For this purpose, we introduce rules that allow for the reasoning about observations. Rules (5)–(7) allow for reasoning about observations of the form \( \text{observed}(l, t) \). The following rules are for observations of the form \( \text{occurred}(a, t) \). For each action \( a \) with executability condition \( h(p_{1}, t), \ldots, h(p_{n}, t) \), we include the set of rules

\[ h(p_{i}, t) \leftarrow \text{occurred}(a, t), \text{not } ab_{i}(t). \quad (i = 1, \ldots, n) \quad (16) \]

This rule says that, if the action \( a \) was observed to occur, then its preconditions must have been satisfied, unless it is an exception. We will denote the set of rules (5)–(7) and (16) as \( \Pi_{R}(KB) \). We define
Definition 2 Given a knowledge base $KB = \langle \Pi_A, \Pi_K, \Pi_O \rangle$, we say that $S$ is an answer set of $KB$ if and only if it is an answer set of the program $\Pi_A \cup \Pi_K \cup \Pi_O \cup \Pi_R(KB)$ under the semantics given in the background section.

Next, we define the notion of a statement.

Definition 3 A statement about $l$ of a language $L$ at the time step $t$ is an expression of the form $\langle l, t \rangle$.

Intuitively, the expression represents a statement by an external source that $l$ holds at the time step $t$. The next definition allows for the evaluation of the truthfulness of a statement against a knowledge base.

Definition 4 Let $KB = \langle \Pi_A, \Pi_K, \Pi_O \rangle$ be a knowledge base over $L$. We say that

- $\text{statement}(l, t)$ is true w.r.t. $KB$, denoted $KB \models +\text{statement}(l, t)$, if for every answer set $S$ of $KB$, $h(l, t) \in S$.
- $\text{statement}(l, t)$ is false w.r.t. $KB$, denoted $KB \models \neg\text{statement}(l, t)$, if for every answer set $S$ of $KB$, $h(l, t) \notin S$.
- $\text{statement}(l, t)$ is unknown w.r.t. $KB$, denoted by $KB \not\models \pm\text{statement}(l, t)$, if $KB \not\models +\text{statement}(l, t)$ and $KB \not\models \neg\text{statement}(l, t)$.

The above definition allows for reasoning about statements against a knowledge base. We consider statements that are true, false, and unknown as truthful, dishonest, and undecided respectively.

Example 1 Let $KB^1_{John} = \langle \Pi_A^1, \Pi_K^1, \Pi_O^1 \rangle$ where

- $\Pi_A^1 = \emptyset$
- $\Pi_K^1$ is the set of rules (4), (9), (11), (13), (14)–(15)
- $\Pi_O^1$ is the set of facts (3), (8), (10), and (12).

We can easily check that

- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 0)$
- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 1)$
- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 2)$
- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 3)$

This reflects correctly our sentiment regarding John’s statement about him not having money given the sequence of observations about John.

Next, we change the example slightly to account for observations about action occurrences.

Example 2 Consider the action of buying an expensive car like a Ferrari (buy_car). We know that a pre-condition of this action is money, i.e., an agent executing this action must have money. Let us know consider the $KB^1_{John} = \langle \Pi_A^1, \Pi_K^1, \Pi_O^1 \rangle$ where

- $\Pi_A^1$ consists of the encoding of the action buy_car (e.g., following the representation described in (Gelfond and Lifschitz 1992)).
- $\Pi_K^1$ is the collection of rules (4), (9), (11), (13), and (14)–(15)
- $\Pi_O^1$ is the set of facts (3), (10), and (12), and the new observation, shown below; instead of (8).

on the observation, shown below, instead of (8).

\begin{equation}
\text{occurred}(\text{buy}_\text{car}, 1)
\end{equation}

Notice that $\Pi_R$ for $KB^1_{John}$ contains the following rule

\begin{equation}
\text{has}_\text{money}(T) \leftarrow \text{occurred}(\text{buy}_\text{car}, T),
\end{equation}

that is an instance of (16). It is easy to see that

- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 0)$
- $KB^1_{John} \models \neg\text{statement}(\neg\text{has}_\text{money}, 1)$
- $KB^1_{John} \models \neg\text{statement}(\neg\text{has}_\text{money}, 2)$
- $KB^1_{John} \models \text{statement}(\neg\text{has}_\text{money}, 3)$

i.e., our conclusions on the statement do not change.

Conclusions and Discussion

We describe a framework based on answer set programming for reasoning about the truthfulness of agents. We show that the task can be achieved by combining the three components, our knowledge about the agent’s actions, our commonsense knowledge about the world, and our observations about the agents. To the best of our knowledge, no such formalization exists.

In the literature, dishonesty and lying have been investigated from different perspective. The focus has been on defining what is a lie (Mahon 2008) or a bullshit (Frankfurt 2005) in a static environment. Our framework takes into considerations changes in the environment that can be the result of the agents’ own actions and/or our knowledge about normal behaviors. So far, we only attempted to evaluate simple statements, i.e., statements about literals. One of our immediate desires is to extend the framework to allow the reasoning about the truthfulness of statements about rules, e.g., “John said that he does not attend college because he does not have money.”

References


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