Learning CP-net Preferences Online from User Queries*

Joshua T. Guerin

The University of Tennessee at Martin 110 Business Administration Building Martin, Tennessee 38238 jguerin@utm.edu

Thomas E. Allen

University of Kentucky
Department of Computer Science
Davis Marksbury Building
329 Rose Street
Lexington, Kentucky 40506-0633
thomas.allen@uky.edu

Judy Goldsmith

University of Kentucky
Department of Computer Science
311 Davis Marksbury Building
329 Rose Street
Lexington, Kentucky 40506-0633
goldsmit@cs.uky.edu

Introduction

CP-nets (Boutilier et al. 1999) offer a compact qualitative representation of human preferences that operate under ceteris paribus ("with all else being equal") semantics. In this paper we present a novel algorithm through which an agent learns the preferences of a user. CP-nets are used to represent such preferences and are learned online through a series of queries generated by the algorithm. Our algorithm builds a CP-net for the user by creating nodes and initializing CPTs, then gradually adding edges and forming more complex CPTs consistent with responses to queries until a confidence parameter is reached. Our algorithm does not always converge to the original CP-net, but our experiments show that it can learn a CP-net that closely tracks with the original for a series of outcome comparison queries. Our work builds upon previous CP-net learning research, particularly that of (Lang and Mengin 2008; 2009) and (Dimopoulos, Michael, and Athienitou 2009). Other CP-net learning algorithms include (Eckhardt and Vojtáš 2009; 2010), (Eckhardt and Vojtáš 2009; 2010), (Koriche and Zanuttini 2009; 2010), and (Liu et al. 2012). Our algorithm differs in that in is guaranteed to produce a CP-net in polynomial time given a constant bound on the number of parents.

Modeling Preferences with CP-nets

By preference, we mean a strict partial order \succ over a set of outcomes \mathcal{O} . Such outcomes can be factored into variables \mathcal{V} with associated (binary) domains $\mathrm{Dom}(\mathcal{V})$: $\mathcal{O} = v_1 \times v_2 \times \cdots \times v_k$. We define o[i] as the projection of outcome o onto variable v_i . Note that the number of outcomes and orderings is exponential in the number of variables. Conditional Preference networks (CP-nets) generally offer a more compact representation.

Definition 1. A CP-net \mathcal{N} is a directed graph. Each node v_i represents a preference over a finite domain. An edge (v_i, v_j) indicates that the preference over v_j depends on v_i . If a node has no incoming edges, the preference involving its variable is not conditioned on other variables. A conditional preference table (CPT) is associated with each node v and

specifies the preference over Dom(v) as a function of the values assigned to its parent nodes Pa(v). A separable CP-net is one with no edges—no variable depends on any other.

To guarantee tractability, we make some simplifying assumptions: 1. Cycles are disallowed. 2. We restrict to binary domains. 3. A maximum bound p is placed on the number of parents a node may have: We conjecture that most human preferences are conditioned on 3–5 nodes and thus feel justified in assuming such a bound.

Algorithm

Our algorithm consists of two phases. First, it constructs a separable CP-net with default CPTs. Next, it successively attempts to refine the model, adding edges and learning more complex CPTs consistent with evidence drawn from the *user queries*. (See LEARN-CP-NET and its subroutine FIND-PARENTS [Alg. 1 and 2]).

Phase 1 constructs a *separable CP-net basis* by asking the user to provide a default preference for each $v_i \in \mathcal{V}$.

Definition 2. Let v_i be a variable in a CP-net with binary domain $Dom(v_i) = \{x_i, y_i\}$. An attribute comparison query is one in which we present the user the values x_i and y_i and ask whether $x_i > y_i$ or $y_i > x_i$.

The result is a CP-net with no edges and only the default values. However, we are *unconfident* that all preferences are unconditional. Here we model *confidence* q as a parameter in our algorithm, defining disjoint sets CONFIDENT and UNCONFIDENT s.t. $v \in \text{CONFIDENT}$ iff we are confident that the preferences over v are conditioned only by its parent variables in the graph of \mathcal{N} .

In the second phase we refine \mathcal{N} by discovering such conditional relationships as may exist between variables by asking the user's preference over pairs of outcomes.

Definition 3. In an outcome comparison query, we provide the user a pair of outcomes, $\{o_1, o_2\} \in \mathcal{O}$. The user responds with $o_1 \succ o_2$, $o_2 \succ o_1$ or $o_1 \sim o_2$, indicating that the user strictly prefers the first outcome to the second, the second to the first, or is indifferent.

Definition 4. A random query is an outcome comparison query in which all values of o_1 and o_2 are selected uniformly randomly from their domains, with the requirement that the query must be relevant to node v_i : that is, $o_1[i] \neq o_2[i]$. A

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Algorithm 1 LEARN-CP-NET(\mathcal{V}, p, q)

```
1: \mathcal{N} \leftarrow \emptyset; comparisons \leftarrow \emptyset
 2: confident \leftarrow \emptyset; unconfident \leftarrow \mathcal{V}
     for v_i \in \mathcal{V} do
         query user: do you prefer x_i \succ y_i or y_i \succ x_i?
 4:
 5:
         v_i.CPT \leftarrow default CPT based on user response
 6:
         insert v_i into \mathcal{N}
 7: end for
 8:
     repeat
 9:
         for r \leftarrow 0 to p do
10:
             for v_i \in unconfident do
                 (P, C) \leftarrow \text{FIND-PARENTS}(v_i, r, q)
11:
12:
                if C \neq FAIL then
13:
                    v_i.\text{CPT} \leftarrow C
14:
                    add edges from all P to v_i
15:
                    move v_i from unconfident to confident
16:
                end if
17:
             end for
18:
         end for
19: until no parents added this iteration
20: return \mathcal{N}
```

Algorithm 2 FIND-PARENTS (v_i, r, q)

```
1: for P \in \{\text{all subsets of } confident \text{ of size } r\} do
        (C, evidCount) \leftarrow CREATE-CPT(v_i, P)
       while (C \neq FAIL) and (evidCount < q) do
3:
 4:
           (o_1, o_2) \leftarrow generate random query for v_i
 5:
           query user: do you prefer outcome o_1 or o_2?
 6:
          add o_1, o_2 to comparisons in specified order
 7:
           (C, evidCount) \leftarrow CREATE-CPT(v_i, P)
 8:
       if C \neq \text{FAIL}, return (C, P), end if
9:
10: end for
11: return (FAIL, ∅)
```

random adaptive query adds the additional requirement that for all $v_i \in \text{CONFIDENT}$, $o_1[j] = o_2[j]$.

Random adaptive queries provide a heuristic that may reduce the search space for a CP-net by not continuing to analyze nodes once they are labeled CONFIDENT.

We search first for nodes that do not need parents. For each node $v \in \mathcal{V}$, we ask a series of outcome comparison queries, then iterate over orderings provided by the user and stored in COMPARISONS. If the user prefers $x_i \succ y_i$ or $y_i \succ$ x_i in all instances, we conclude that the preferences over v_i are unconditional and move it from UNCONFIDENT to CONFIDENT. If we have not accumulated enough evidence, we continue querying the user. While UNCONFIDENT $\neq \emptyset$, we continue trying to refine our model with new conditional relationships, represented as edges and more complex CPTs. For each unconfident node, we iterate over potential sets of parent nodes of increasing size up to p. If, in an iteration, we fail to add parents for any nodes, we stop. Our algorithm will always output a CP-net, possibly with some CPTs in their default state from Phase 1; however, this rarely occurred in our tests, and only in overtrained CP-nets of minimal size.

For a given target node and set of possible parents, we construct a 2-SAT instance such that (1) a satisfying assignment tells us that the target node's values are consistent with

the given set of parents and (2) the assignment to variables gives us the entries of the target node's CPT. Our method for this closely follows (Dimopoulos, Michael, and Athienitou 2009), to which the reader is referred for specifics.

Analysis and Experiments

Theorem 1. LEARN-CP-NET is resolute—i.e., it is guaranteed to output a consistent CP-net \mathcal{N} —and runs in time polynomial in n^p and q in the worst case. (Proof omitted.)

We generated random CP-nets for given n and p, and used them to generate responses to the queries for our learning algorithms. We looked at computation time (as a function of n, p, and q), and the accuracy of the learned CP-net, i.e., on how many possible comparison queries do the generated and learned CP-nets agree. The tables below show the metrics of the learned CP-net \mathcal{N}_L compared with the training model \mathcal{N}_T over a series of experiments.¹ We set p=5; since we used $\delta = n$, most nodes didn't have p parents. However, in trials with $\delta = cn$ for c = 2, 3, ..., we saw very similar graphs. The metrics shown are averages over 10 trials. Table 1 shows that agreement was generally 75-90%+ with the proper q. As shown, disagreement between models was rare, but as n increases, the learned model is more likely to be indecisive about preferences on which the training model decides. Increasing q sometimes has an adverse effect on the agreement the models; if q is too high, the model can be overtrained. We also found that for some q values, computational time did not grow monotonically. When we generate queries to learn the CPT for v_i , those queries may be relevant to other nodes in UNCONFIDENT. It may be that, when we come to v_i , we already have q many relevant comparisons.

Table 1: Agreement of \mathcal{N}_L with \mathcal{N}_T

q	n=3	n=5	n = 7	n = 10
4	0.9964	0.7996	0.6086	0.5118
6	0.9964	0.9221	0.7653	0.7677
8	1.0000	0.9667	0.8918	0.6990
10	0.9964	0.9735	0.8583	0.6760
12	0.9964	0.9621	0.9171	0.6597
14	1.0000	0.9816	0.8512	0.5518
16	0.9929	0.9434	0.9062	0.5539
18	1.0000	0.9646	0.9237	0.4367
20	0.9964	0.9731	0.8395	0.5886

Table 2: Disagreement of \mathcal{N}_L with \mathcal{N}_T

q	n=3	n=5	n=7	n = 10
4	0.0036	0.0696	0.0880	0.0426
6	0.0036	0.0430	0.0545	0.0389
8	0	0.0238	0.0375	0.0280
10	0.0036	0.0216	0.0257	0.0254
12	0.0036	0.0174	0.0188	0.0150
14	0	0.0184	0.0259	0.0118
16	0.0071	0.0309	0.0250	0.0127
18	0	0.0228	0.0186	0.0095
20	0.0036	0.0269	0.0236	0.0110

¹Due to space constraints, indecision results have been omitted; these values can be computed from Tables 1–2.

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